

AMENDMENTS TO THE CLAIMS:

This listing of claims will replace all prior versions, and listings, of claims in the application:

LISTING OF CLAIMS:

1. (Currently Amended) A cryptographic method during which use is made of a random number generator producing random numbers S_i of size N fixed between 0 and $W-1$, in order to produce a random number R between 0 and a predefined limiter K , characterised ~~in that~~ wherein:

E31: a random variable S_i between 0 and $W-1$ is produced,

E32: if the random variable S_i is strictly less than a coefficient K_i of the limiter K in base W , then the coefficient R_i of rank i of the random number R is equal to the random variable S_i and then, for any rank J less than i , a random variable S_j between 0 and $W-1$ is produced and $R_j = S_j$,

E33: otherwise, if the said random variable is greater than the coefficient K_i of rank i of the limiter K in base W , then the said coefficient R_i is determined from the random variable S_i of rank i according to a predetermined function, and then the coefficient R_{i-1} is determined for the random number R of rank $i-1$ that is immediately lower by repeating steps E31 to E33.

2. (Currently Amended) A method according to claim ~~[[2]]~~ 1, during which the following steps are performed:

E1: the limiter K is decomposed in base $(W^{p-1}, W^{p-2}, \dots, W^0)$ in the form

$$K = \sum_{i=0}^{p-1} K_i * W^i$$
, i being a loop index, K_i being a coefficient of the limiter K of rank i between 0 and $W-1$ and p being the degree of the limiter K ,

E2: a Boolean variable f is initialised to TRUE,

E3: the following operations are performed, in a loop indexed by i , i being an integer varying between $p-1$ and 0:

E31: a random variable S_i between 0 and ~~[[W0-1]]~~ $W-1$ is produced,

E32: if the random variable S_i is strictly less than the coefficient K_i of rank i , then the Boolean variable f is set to FALSE,

E33_1: if the random variable S_i is strictly greater than the coefficient K_i of rank i and the Boolean variable f is TRUE, then the coefficient R_i of rank i is determined from the random variable S_i of rank i according to a predefined function,

E33_2: otherwise $R_i = S_i$

E34: the loop ~~indexed~~ index i is decremented,

E4: the random number R is determined by recombination of the random coefficients

$$R = \sum_{i=0}^{p-1} R_i * W^i$$

R_i in base W according to the equation:

3. (Original) A method according to claim 2, during which, in order to determine the coefficient R_i of rank i from the random variable S_i of rank i (steps E33_1 and E33_2), the following substeps are performed:

E33_11: if the random variable S_i is strictly greater than the coefficient K_i of the limiter K , then a new random variable S_i is produced,

E33_12: step E33_11 is repeated until the random variable S_i is less than the coefficient K_i of the limiter K , and then the coefficient R_i is equalised to the random variable S_i .

4. (Currently Amended) A method according to claim 2, during which the coefficient R_i of rank i is chosen (steps E33-1 and E33_2) equal to the part of the random variable S_i , the part less than the coefficient K_i , the said part corresponding to a limited number of bits of the variable S_i .

5. (Original) A method according to claim 2, during which, in order to determine the coefficient R_i of rank i from the random variable S_i of rank i (step E33), the random variable S_i is reduced modulo K_i+1 , the result of the reduction being the coefficient sought.

6. (Currently Amended) A method according to ~~one of claims 1 to 5~~ claim 2, during which, in order to determine the coefficient R_i of rank i from the random variable S_i of rank i (step E33), steps E1 to E4 are executed using a base $(\beta^{q-1}, \dots, \beta^0)$ as the calculation base, β being an integer strictly less than W and q being the degree of k in case β .

7. (Currently Amended) A method according to claim 6, in which step E33 is broken down into the following substeps:

E33_41: the coefficient K_i of rank i of the limiter K in base $(\beta^{q-1}, \dots, \beta^0)$ in the form

$$K_i = \sum_{j=0}^{q-1} (K_i)_j * \beta^j$$

, j being a loop index, $(K_i)_j$ being a number between 0 and $\beta-1$ and q being a degree of the coefficient K_i , is decomposed,

E33_42: a second Boolean variable g is initialised to TRUE,

E33_43: the following operations are performed, in a loop indexed by j varying between $q-1$ and 0:

E33_431: a random variable $(S_i)_j$ between 0 and $\beta-1$ is produced,

E33_432: if the random variable $(S_i)_j$ is strictly less than the coefficient $(K_i)_j$, then the second Boolean variable g is set to FALSE,

E33_4331: if the random variable $(S_i)_j$ is strictly greater than the coefficient $(K_i)_j$ and the second Boolean variable g is TRUE, then a coefficient $(R_i)_j$ is determined from the random variable $(S_i)_j$ according to a predefined function,

E33_4332: otherwise, $(R_i)_j = (S_i)_j$

E33_434: the loop ~~indexed~~ index j is decremented,

E33_44: the random number R_i is determined by recombination of the random

$$R_i = \sum_{j=0}^{q-1} (R_i)_j * \beta^j$$

coefficients $(R_i)_j$ in base β according to the equation:

8. (Currently Amended) An electronic component comprising a generator of random numbers of size N , calculation circuits performing in particular a comparison, a truncation and/or a modular reduction on numbers of no more than N bits, and a means of controlling the random number generator and calculation circuits, the said control means being adapted for implementing a method according to ~~one of claims 1 to 7~~ claim 1.

9. (Currently Amended) A chip card comprising an electronic component according to ~~the preceding claim~~ claim 1.

10. (New) A method according to claim 3, during which, in order to determine the coefficient R_i of rank i from the random variable S_i of rank i (step E33), steps E1 to E4 are executed using a base $(\beta^{q-1}, \dots, \beta^0)$ as the calculation base, β being an integer strictly less than W and q being the degree of k in case β .

11. (New) A method according to claim 4, during which, in order to determine the coefficient R_i of rank i from the random variable S_i of rank i (step E33), steps E1 to E4 are executed using a base $(\beta^{q-1}, \dots, \beta^0)$ as the calculation base, β being an integer strictly less than W and q being the degree of k in case β .

12. (New) A method according to claim 5, during which, in order to determine the coefficient R_i of rank i from the random variable S_i of rank i (step E33), steps E1 to E4 are executed using a base $(\beta^{q-1}, \dots, \beta^0)$ as the calculation base, β being an integer strictly less than W and q being the degree of k in case β .

13. (New) An electronic component comprising a generator of random numbers of size N , calculation circuits performing in particular a comparison, a truncation and/or a modular reduction on numbers of no more than N bits, and a means of controlling the random number generator and calculation circuits, the said control means being adapted for implementing a method according to claim 2.

14. (New) An electronic component comprising a generator of random numbers of size N , calculation circuits performing in particular a comparison, a truncation and/or a modular reduction on numbers of no more than N bits, and a means of controlling the random number generator and calculation circuits, the said control means being adapted for implementing a method according to claim 3.

15. (New) An electronic component comprising a generator of random numbers of size N , calculation circuits performing in particular a comparison, a truncation and/or a modular reduction on numbers of no more than N bits, and a means of controlling the random number generator and calculation circuits, the said control means being adapted for implementing a method according to claim 4.

16. (New) An electronic component comprising a generator of random numbers of size N , calculation circuits performing in particular a comparison, a truncation and/or a modular reduction on numbers of no more than N bits, and a means of controlling the random number generator and calculation circuits, the said control means being adapted for implementing a method according to claim 5.

17. (New) An electronic component comprising a generator of random numbers of size N, calculation circuits performing in particular a comparison, a truncation and/or a modular reduction on numbers of no more than N bits, and a means of controlling the random number generator and calculation circuits, the said control means being adapted for implementing a method according to claim 6.

18. (New) An electronic component comprising a generator of random numbers of size N, calculation circuits performing in particular a comparison, a truncation and/or a modular reduction on numbers of no more than N bits, and a means of controlling the random number generator and calculation circuits, the said control means being adapted for implementing a method according to claim 7.